American University of Beirut


# Discrete Structures Topic 1 - Logic: Propositional Logic (Ch 1.1)* 

CMPS 211 - Fall 2017 - American University of Beirut

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## Why logic?

- Logic is a set of principles that can be used to reason about (mathematical) statements
- For instance, let's say we want to formally express and reason about the following statement:
* "For every positive integer $n$, the sum of the positive integers not exceeding $n$ is $n(n+1) / 2^{\prime \prime}$
- We can formally express the above statement using logic
- We can also prove the above statement or argue whether it is true or false using logic


## Why logic? (cont.)

- Logic has numerous applications to Computer Science
- Used in the design of computer circuits
- Used in the construction of computer programs
- Used to verify the correctness of programs
- Used to ensure the security of a system
- Used heavily in artificial intelligence


## Propositional Logic

## Propositions

- A proposition is a declarative statement that is either true or false
- Examples of propositions:
a) The Moon is made of white cheese
b) Toronto is the capital of Canada
c) A week has more days than a month
d) $1+0=1$
e) $0+0=2$
- Examples that are not proposition
a) Sit down! - command
b) What time is it? - question
c) $1+2$-expressions with a non-true/false value
d) This statement is false - a paradox


## Atomic propositions

- We use letters ( $\mathrm{p}, \mathrm{q}, \mathrm{r}, \mathrm{s}, \ldots$ ) to denote atomic propositions
- Also called propositional variables
- Similar to $x, y, z, \ldots$ for numerical variables
- For example, let p be the proposition that the earth is round and $q$ be the proposition that the moon is flat
- These represent a single statement that cannot be "decomposed"


## Compound propositions

- Compound propositions are built up from atomic propositions by the use of Boolean connectives. Also called propositional formulae.
- Propositional Logic is the logic of propositional formulae and their meaning
- First developed by the Greek philosopher Aristotle more than 2300 years ago
- George Boole introduced Boolean Algebra in 1854


## Propositional Connectives

## Propositional Operators/Connectives

- An operator or connective combines one or more operand expressions into a larger expression (e.g., "+" in numeric expressions)
- Unary operators take 1 operand (e.g., -3)
- binary operators take 2 operands (e.g., $3 \times 4$ )
- Propositional or Boolean operators operate on propositions (or their truth values) instead of numbers
- There are six main operators
- Negation $\neg$
- Conjunction $\wedge$
- Disjunction V
- The Exclusive Or $\oplus$
- Implication $\rightarrow$
- Biconditional $\leftrightarrow$


## Connectives: Negation

- The negation of a proposition $p$ is denoted by $\neg p$
- In an English statement, we express $\neg p$ as follows: "It's not the case that $p$ "
- $\neg$ p is true if $p$ is false and is false if $p$ is true
- Example:
- If p denotes "I am at home",
- then $\neg \mathrm{p}$ denotes "It is not the case that I am at home" or more simply "I am not at home"
- The negation of a proposition $p$ has this truth table

| $p$ | $\neg p$ |
| :--- | :--- |
| T | F |
| F | T |

## Connectives: Conjunction

- The conjunction of propositions $p$ and $q$ is denoted by $p \wedge q$
- It is the proposition " $p$ and $q$ "
- $p \wedge q$ is true when both $p$ and $q$ are true and is false otherwise
- Example:
- If $p$ denotes "I am at home", and
- $q$ denotes "It is raining"
- then $p \wedge q$ denotes "I am at home and it is raining"


## Conjunction (cont.)

- The conjunction of propositions $p$ and $q$ has this truth table

| $p$ | $q$ | $p \wedge q$ |
| :--- | :--- | :--- |
| T | T | T |
| T | F | F |
| F | T | F |
| F | F | F |

## Connectives: Disjunction

- The disjunction of propositions $p$ and $q$ is denoted by $p \vee q$
- It is the proposition " $p$ or $q$ "
- $p \vee q$ is false when both $p$ and $q$ are false and is true otherwise
- Example:
- If $p$ denotes "I am at home", and
- $q$ denotes "It is raining"
- then $p \vee q$ denotes "I am at home or it is raining"


## Disjunction (cont.)

- The disjunction of propositions $p$ and $q$ has this truth table

| $p$ | $q$ | $p \vee q$ |
| :--- | :--- | :--- |
| T | T | T |
| T | F | T |
| F | T | T |
| F | F | F |

## Connectives: The Exclusive Or

- The exclusive or of $p$ and $q$ is denoted by $p \oplus q$
- We say $p$ XOR $q$
- $p \oplus q$ is true when exactly either $p$ or $q$ is true and is false otherwise
- Example:
- If $p$ denotes "I am at home", and
- $q$ denotes "It is raining"
- then $p \oplus q$ denotes :"either I am at home or it is raining" but not both


## The Exclusive Or (cont.)

- Note that English "or" can be ambiguous regarding the "both" case!
- "Pat is a singer or

Pat is a writer" V

- "Pat is a man or

Pat is a woman"$\oplus$

- Need context to understand the meaning!
- For this class, assume "or" means inclusive


## The Exclusive Or (cont.)

- The exclusive or of propositions $p$ and $q$ has this truth table

| $p$ | $q$ | $p \oplus q$ |
| :--- | :--- | :--- |
| T | T | F |
| T | F | T |
| F | T | T |
| F | F | F |

## Connectives: Implication

- The conditional statement or implication $p \rightarrow q$ is the proposition "if $p$, then $q$ "
- $p \rightarrow q$ is false when $p$ is true and $q$ is false, and is true otherwise
- $p$ is called the hypothesis (or antecedent or premise) and $q$ is called the conclusion (or consequence)
- Example:
- If $p$ denotes "I am at home", and
- $q$ denotes "It is raining"
- then $p \rightarrow q$ denotes "If I am at home then it is raining"


## Implication (cont.)

- The implication or conditional statement $p \rightarrow q$ has this truth table

| $p$ | $q$ | $p \rightarrow q$ |
| :--- | :--- | :--- |
| T | T | T |
| T | F | F |
| F | T | T |
| F | F | T |

## Understanding Implication

- In $p \rightarrow q$ there does not need to be any connection between the hypothesis and the conclusion
- These implications are perfectly fine, but would not be used in ordinary English
- "If the moon is made of green cheese, then I have more money than Bill Gates "
- "If $1+1=3$, then pink elephants can fly"


## Examples

* "If you get $100 \%$ on the final, then you will get an A"
- Interpretation:
- If you manage to get a $100 \%$ on the final, then you would expect to receive an A
- If you do not get a $100 \%$, you may or may not receive an A, depending on other factors (such as...)
- However, if you do get 100\%, but the professor does not give you an A, you will feel cheated


## More Examples

" "When I got elected, I will lower the taxes"

- Interpretation:
- If the politician is elected, voters would expect the taxes to get lower
- If the politician is not elected, then the voters have no expectations regarding the taxes, it might get lower, higher, or stay the same.
- It is only when the politician is elected and the taxes are not lower, the voters would say that the politician has broken his campaign pledge


## Different Ways of Expressing $p \rightarrow q$

- if $p$, then $q$
- if $p, q$
- $p$ implies $q$
- $p$ only if $q$
- a necessary condition for $p$ is $q$
- $p$ is sufficient for $q$
- $q$ if $p$
- $q$ whenever $p$
- $q$ when $p$
- $q$ follows from $p$
- q unless $\neg p$
- $q$ is necessary for $p$
- a sufficient condition for $q$ is $p$


## Sufficient

* "if $p$ then $q$ " expresses the same thing as " $p$ is sufficient for $q$ "
- This is basically saying that $p$ holding is sufficient for concluding that $q$ will also hold
- Example:
, "if $p$ then $q$ ": If Maria learns discrete mathematics, then she will find a good job

| $p$ | $q$ | $p \rightarrow q$ |
| :---: | :---: | :---: |
| F | F | T |
| F | T | T |
| $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |
| T | T | T |

" " $p$ is sufficient for $q$ ": Learning discrete mathematics is sufficient for Maria to find a good job

- If Maria doesn't learn discrete mathematics, she might find a good job or not, however if she does learn discrete mathematics, she will find a good job for sure


## Necessary

- "if $p$ then $q$ " expresses the same thing as " $q$ is necessary for $p$ "
- This is basically saying that q holding is a necessary conclusion for a holding premise $p$
- Example:
- 'if $p$ then $q$ ": If you can access the Internet, then

| $p$ | $q$ | $p \rightarrow q$ |
| :---: | :---: | :---: |
| F | F | T |
| F | T | T |
| $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |
| T | T | T | you have paid a subscription fee

- " $q$ is necessary for $p$ ": Paying a subscription fee is necessary for being able to access the Internet
- If you can access the Internet, then you must have paid a subscription fee, however if you cannot access the Internet, you may have paid a subscription fee or not


## Only If

* "if $p$ then $q$ " expresses the same thing as " $p$ only if $q^{\prime \prime}$
- This is basically saying that $p$ cannot be true when $q$ is not true
- Example:
* "if $p$ then $q$ ": If you can access the Internet,

| $p$ | $q$ | $p \rightarrow q$ |
| :---: | :---: | :---: |
| F | F | T |
| F | T | T |
| $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |
| T | T | T | then you have paid a subscription fee"

- " $p$ only if $q$ ": You can access the Internet only if you pay a subscription fee
- If you haven't paid a subscription fee, you cannot access the Internet, however if you did pay a subscription fee, you might be able to access the Internet or not


## Unless

- "if $p$ then $q$ " expresses the same thing as " $q$ unless $\neg p^{\prime \prime}$
- This is basically saying that $q$ will be true except if $p$ is false
- Example:
- "if $p$ then $q$ ": If Maria learns discrete

| $p$ | $q$ | $p \rightarrow q$ |
| :---: | :---: | :---: |
| F | F | T |
| F | T | T |
| $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |
| T | T | T |

Mathematics, then she will find a good job

- "q unless $\neg$ p": Maria will find a good job unless she does not learn discrete mathematics
- We surely know Maria will get a good job, unless she doesn't learn discrete Mathematics - in that case, we don't know whether she'll find a good job or not


## Converse, Contrapositive, and Inverse

- From $p \rightarrow q$, we can form new conditional statements
- $q \rightarrow p \quad$ is the converse of $p \rightarrow q$
- $\neg q \rightarrow \neg p$ is the contrapositive of $p \rightarrow q$
- $\neg p \rightarrow \neg q$ is the inverse of $p \rightarrow q$
- Out of these three conditional statements formed from $p \rightarrow q$, only the contrapositive always has the same truth value as $p \rightarrow q$
- The other two might seem to have a similar meaning, but actually translate to different premises and conclusions


## Example

Let $p$ denote "I am elected president", and $q$ denote "I will make healthcare free"

- Original $p \rightarrow q$ : "If I am elected president, I will make healthcare free"
- Converse $q \rightarrow p$ : "If the healthcare was made free, then I am elected president"
- Contrapositive $\neg q \rightarrow \neg p$ : "If healthcare wasn't made free, then I haven't been elected president"
- Inverse $\neg p \rightarrow \neg q$ : "If I wasn't elected president, then the health care won't be made free"


## Example Explained

- The contrapositive clearly means the same thing as the original statement
- I made a pledge that healthcare will be made free when I get elected to presidency
- Consequently, if the healthcare wasn't made free, then surely I haven't been elected for presidency
- The converse, on the other hand, means something else
- It says that observing healthcare made for free will lead to concluding that I have been elected for presidency, which isn't what the original statement says
- The original statement keeps space for the case when I don't get elected for presidency, yet someone else makes healthcare for free
- The same logic applies for the inverse


## Exercise

- Find the converse, inverse, and contrapositive of "It raining is a sufficient condition for me not going to town"
- Solution:
- Original : If it is raining, then I won't go to town
- Converse: If I do not go to town, then it is raining
- Inverse: If it is not raining, then I will go to town
- Contrapositive: If I go to town, then it is not raining


## Another Exercise

- Find the converse, inverse, and contrapositive of "The home team wins whenever it is raining"
- Solution:
- Original: if it is raining, then the home team wins
- Converse: if the home team wins, then it is raining
- Inverse: if it is not raining, the home team won't win
- Contrapositive: If the home team doesn't win, then it is not raining


## Connectives: Biconditional

- The biconditional proposition $p \leftrightarrow q$ is the proposition " $p$ if and only if $q$ "
- The biconditional $p \leftrightarrow q$ is true when p and q have the same truth values and is false otherwise
- Also called bi-implications
- Example:
- If $p$ denotes "buying a ticket",
- and $q$ denotes "can take a flight"
- then $p \leftrightarrow q$ denotes "I can take a flight if and only if I buy a ticket"
- This statement is true when you buy a ticket and can take a flight, Or when you don't buy a ticket and can't take a flight
- It is false when you buy a ticket and can't take a flight (like when the airlines bumps you)
- Or when you don't buy a ticket and can take a flight (like when you win a ticket)


## Biconditional

- The biconditional proposition $p \leftrightarrow q$ has this truth table

| $p$ | $q$ | $p \leftrightarrow q$ |
| :--- | :--- | :--- |
| T | T | T |
| T | F | F |
| F | T | F |
| F | F | T |

## Expressing the Biconditional

- Some alternative ways " $p$ if and only if $q$ " is expressed in English
- $p$ iff $q$
- if $p$ then $q$, and conversely
- $p$ is necessary and sufficient for $q$
- Example:
- You will be always informed if and only if you read a newspaper everyday
- If you read a newspaper everyday, then you will be informed, and conversely
- To be always informed, it is necessary and sufficient to read a newspaper everyday


## Implicit Use Of Biconditionals

- Human natural language is not explicit on biconditionals most of the times
- People use "if, then" and "only if" constructs and mean "if and only if" implicitly
- For example, consider the statement in English: "If you finish your meal, then you can have dessert"
- What is really meant is "You can have dessert if and only if you finish your meal"
- Here, we will always distinguish between the conditional statement $p \rightarrow q$ and the biconditional statement $p \leftrightarrow q$


## Precedence of Logical Operators

- When multiple logical operators are used in the same compound proposition, we follow the following precedence order to understand it:
- Note that all operators are right associative (conventionally)
- Parentheses can also be used to disambiguate
- Example:

Operator Precedence
, $p \vee q \rightarrow \neg r$ is equivalent to $(p \vee q) \rightarrow \neg r$

- If the intended meaning is $p \vee(q \rightarrow \neg r)$, then parentheses must be used

| $\boldsymbol{\neg}$ | 1 |
| :---: | :---: |
| $\boldsymbol{\wedge}$ | 2 |
| $\mathbf{v}$ | 3 |
| $\rightarrow$ | 4 |
| $\leftrightarrow$ | 5 |

## Another Example

- Parse the following propositional statement Op. Prec.

$$
\mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s}
$$

| $\neg$ | 1 |
| :---: | :---: |
| $\wedge$ | 2 |
| $\mathbf{v}$ | 3 |
| $\rightarrow$ | 4 |
| $\leftrightarrow$ | 5 |

- Note that all operators are right associative
- Parentheses can also be used to disambiguate


## Another Example

- Parse the following propositional statement Op. Prec.

$$
\begin{gathered}
\mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
\mathrm{p} \wedge(\neg \mathrm{q}) \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s}
\end{gathered}
$$

| $\rightarrow$ | 4 |
| :--- | :--- |
| $\leftrightarrow$ | 5 |

- Note that all operators are right associative
- Parentheses can also be used to disambiguate


## Another Example

- Parse the following propositional statement Op. Prec.

$$
\begin{gathered}
\mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
\mathrm{p} \wedge(\neg \mathrm{q}) \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s}
\end{gathered}
$$

| $\rightarrow$ | 4 |
| :--- | :--- |
| $\leftrightarrow$ | 5 |

- Note that all operators are right associative
- Parentheses can also be used to disambiguate


## Another Example

- Parse the following propositional statement


## Op. Prec.

$$
\begin{aligned}
\mathrm{p} \wedge \neg \mathrm{q} & \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
\mathrm{p} \wedge(\neg \mathrm{q}) & \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) & \rightarrow(\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) & \rightarrow((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}) \leftrightarrow(\mathrm{q} \vee \mathrm{~s})
\end{aligned}
$$

| Op. |
| :--- |
| $-\quad 1$ |1

$\wedge \quad 2$
v 3
$\begin{array}{ll}\rightarrow & 4 \\ \leftrightarrow & 5\end{array}$

- Note that all operators are right associative
- Parentheses can also be used to disambiguate


## Another Example

- Parse the following propositional statement


## Op. <br> Prec.

$$
\begin{aligned}
& \mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
& \mathrm{p} \wedge(\neg \mathrm{q}) \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
& (\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
& (p \wedge(\neg q)) \rightarrow((p \wedge r) \vee s) \rightarrow(\neg t) \leftrightarrow(q \vee s) \\
& (p \wedge(\neg q)) \rightarrow(((p \wedge r) \vee s) \rightarrow(\neg t)) \leftrightarrow(q \vee s)
\end{aligned}
$$

- Note that all operators are right associative
- Parentheses can also be used to disambiguate


## Another Example

- Parse the following propositional statement


## Op. <br> Prec.

$$
\begin{gathered}
\mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
\mathrm{p} \wedge(\neg \mathrm{q}) \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}) \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t})) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}) \\
((\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}))) \leftrightarrow(\mathrm{q} \vee \mathrm{~s})
\end{gathered}
$$

| $\boldsymbol{\jmath}$ | 1 |
| :---: | :---: |
| $\boldsymbol{\wedge}$ | 2 |
| $\mathbf{v}$ | 3 |
| $\rightarrow$ | 4 |
| $\rightarrow$ | 5 |

## Another Example

- Parse the following propositional statement


## Op. <br> Prec.

$$
\begin{gathered}
\mathrm{p} \wedge \neg \mathrm{q} \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow \neg \mathrm{t} \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
\mathrm{p} \wedge(\neg \mathrm{q}) \rightarrow \mathrm{p} \wedge \mathrm{r} \vee \mathrm{~s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s} \rightarrow(\neg \mathrm{t}) \leftrightarrow \mathrm{q} \vee \mathrm{~s} \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}) \\
(\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t})) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}) \\
((\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow(((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}))) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}) \\
((\mathrm{p} \wedge(\neg \mathrm{q})) \rightarrow((\mathrm{p} \wedge \mathrm{r}) \vee \mathrm{s}) \rightarrow(\neg \mathrm{t}))) \leftrightarrow(\mathrm{q} \vee \mathrm{~s}))
\end{gathered}
$$

- Note that all operators are right associative
- Parentheses can also be used to disambiguate

Truth Tables

## Truth Tables For Compound Propositions

- Rows
- Need a row for every possible combination of values for the atomic propositions
- Columns
- Need a column for the compound proposition (usually at far right)
- Need a column for the truth value of each expression that occurs in the compound proposition as it is built up
- This includes the atomic propositions


## Example Truth Table

Construct a truth table for $p \vee q \rightarrow \neg r$

| p | q | r | $\neg \mathrm{r}$ | $\mathrm{p} \vee \mathrm{q}$ | $\mathrm{p} \vee \mathrm{q} \rightarrow \neg \mathrm{r}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| T | T | T | F | T | F |
| T | T | F | T | T | T |
| T | F | T | F | T | F |
| T | F | F | T | T | T |
| F | T | T | F | T | F |
| F | T | F | T | T | T |
| F | F | T | F | F | T |
| F | F | F | T | F | T |

## Equivalent Propositions

- Two propositions are equivalent if they always have the same truth value
- Example:
- Show using a truth table that the implication is equivalent to the contrapositive: $(\mathrm{p} \rightarrow \mathrm{q}) \Leftrightarrow(\neg \mathrm{q} \rightarrow \neg \mathrm{p})$
- Solution:

| $p$ | $q$ | $\neg p$ | $\neg q$ | $p \rightarrow q$ | $\neg q \rightarrow \neg p$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T | T | F | F | T | T |
| T | F | F | T | F | F |
| F | T | T | F | T | T |
| F | F | T | T | T | T |

## Using a Truth Table to Show Non-Equivalence

## - Example:

- Show using truth tables that neither the converse nor inverse of an implication are equivalent to the implication
- Solution:

| $p$ | $q$ | $\neg p$ | $\neg q$ | $p \rightarrow q$ | $\neg p \rightarrow \neg q$ | $q \rightarrow p$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| T | T | F | F | T | T | T |
| T | F | F | T | F | T | T |
| F | T | T | F | T | F | F |
| F | F | T | T | T | T | T |

## Biconditional Truth Table

- $p \leftrightarrow q$ means that $p$ and $q$ have the same truth value
- Note this truth table is the exact opposite of $\oplus$ 's!

Thus, $p \leftrightarrow q$ means $\neg(p \oplus q)$

- $p \leftrightarrow q$ does not imply that $p$ and $q$ are true, or that either of them causes the other, or that they have a common cause

| p | q | $\mathrm{p} \oplus \mathrm{q}$ | $\neg(\mathrm{p} \oplus \mathrm{q})$ | $\mathrm{p} \leftrightarrow \mathrm{q}$ |
| :--- | :--- | :--- | :--- | :--- |
| T | T | F | T | T |
| T | F | T | F | F |
| F | T | T | F | F |
| F | F | F | T | T |

## Exercise

- How many rows are there in a truth table with $n$ propositional variables?
- Solution: $2^{\text {n }}$
- Note that this means that with n propositional variables, we can construct $2^{\text {n }}$ distinct (i.e., not equivalent) propositions


## Semantics of propositional formulae

- A proposition is valid iff it evaluates to true in every row of its truth table
- A proposition is satisfiable iff it evaluates to true in some row of its truth table
- A proposition is valid iff its negation is not satisfiable
- A proposition is a contingency iff it is satisfiable and its negation is satisfiable
- A proposition is a contradiction iff it evalutes to false in every row of its truth table, i.e., iff its negation is valid


## Any Questions?


[^0]:    * Extracted from Discrete Mathematics and It's Applications book slides

