

1. (total 12 pts) Let $A = \{1, 2, 3\}$ and $B = \{1, 2, 5, 10\}$.
 - a) (2 pts) Find $A \cup B$.
 - b) (2 pts) Find $A \cap \overline{B}$. (The answer does not depend on the choice of universal set.)
 - c) (4 pts) Write the following sentence using symbols, and determine whether it is true or false (briefly say why): “Every element of A is a divisor of some element of B ”
 - d) (4 pts) Write the following sentence using symbols, and determine whether it is true or false (briefly say why): “Some element of A is a divisor of every element of B ”

Answer. (a) $A \cup B = \{x \mid (x \in A) \vee (x \in B)\} = \{1, 2, 3, 5, 10\}$.

(b) $A \cap \overline{B} = \{x \mid (x \in A) \wedge (x \notin B)\} = \{3\}$.

(c) $\forall x \in A, \exists y \in B$ s.t. $x|y$. (Alternatively: $\forall x \in A, \exists y \in B, \exists n \in \mathbf{Z}$ s.t. $y = nx$. This statement is false since $3 \in A$ and $\neg(3|1) \wedge \neg(3|2) \wedge \neg(3|5) \wedge \neg(3|10)$.)

(d) $\exists x \in A$ s.t. $\forall y \in B, x|y$.

This statement is true since $1 \in A$ and $(1|1) \wedge (1|2) \wedge (1|5) \wedge (1|10)$.

2. (total 12 pts) Let $f : \mathbf{Z} \rightarrow \mathbf{Z}$ be the function defined by $f(n) = \left\lceil \frac{n}{3} \right\rceil$.
 - a) (2 pts) Find the image set $f(\{1, 3, 4, 100\})$.
 - b) (2 pts) Find the inverse image set $f^{-1}(\{-1, 10\})$.
 - c) (4 pts) Is f an injection? Why or why not (give a brief explanation)?
 - d) (4 pts) Is f a surjection? Why or why not (give a brief explanation)?

Answer. (a) $f(1) = 1, f(3) = 1, f(4) = 2$ and $f(100) = 34$. Hence $f(\{1, 3, 4, 100\}) = \{1, 2, 34\}$.

(b) $f^{-1}(\{-1, 10\}) = \{n \in \mathbf{Z} \mid \left\lceil \frac{n}{3} \right\rceil \in \{-1, 10\}\}$. So for $n \in \mathbf{Z}$, we have

$$\begin{aligned} n \in f^{-1}(\{-1, 10\}) &\leftrightarrow \left(\left\lceil \frac{n}{3} \right\rceil = -1\right) \vee \left(\left\lceil \frac{n}{3} \right\rceil = 10\right) \\ &\leftrightarrow \left(-2 < \frac{n}{3} \leq -1\right) \vee \left(9 < \frac{n}{3} \leq 10\right) \leftrightarrow (-6 < n \leq -3) \vee (27 < n \leq 30). \end{aligned}$$

This means that $f^{-1}(\{-1, 10\}) = \{-5, -4, -3, 28, 29, 30\}$.

(c) f is not an injection since $f(1) = f(2) = 1$ even though $1 \neq 2$. In other words, there are distinct elements of \mathbf{Z} having the same image by f .

(d) f is a surjection since for all $n \in \mathbf{Z}$, we can find a preimage $k = 3n$ for which $f(3n) = n$. Hence every integer n has at least one preimage. (Note that the other preimages of n are $3n - 1$ and $3n - 2$.)

3. (total 12 pts) Show that $p \rightarrow (q \rightarrow r)$ is logically equivalent to $(p \wedge q) \rightarrow r$ in two different ways, (a) and (b+c):
 - a) (4 pts) Use a truth table.
 - b) (4 pts) Give a direct proof that if $p \rightarrow (q \rightarrow r)$, then $(p \wedge q) \rightarrow r$.
 - c) (4 pts) Give a direct proof of the converse of the statement in (b).

Answer. (a) Since the fifth and seventh columns below are identical, we obtain the desired logical equivalence.

p	q	r	$q \rightarrow r$	$p \rightarrow (q \rightarrow r)$	$p \wedge q$	$(p \wedge q) \rightarrow r$
T	T	T	T	T	T	T
T	T	F	F	F	T	F
T	F	T	T	T	F	T
T	F	F	T	T	F	T

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F	T	T	T	T	F	T
F	T	F	F	T	F	T
F	F	T	T	T	F	T
F	F	F	T	T	F	T

(b) We shall give a fairly formal proof, to emphasize the structure of the argument (e.g., where we make an assumption, where we make a subassumption, and so forth). Note how we use indentation to show where an internal argument begins and ends.

- [1] Assumption: $p \rightarrow (q \rightarrow r)$
- [2] Sub-assumption: $p \wedge q$
- [3] therefore p (from [2])
- [4] and q (also from [2])
- [5] therefore $q \rightarrow r$ (combining [3] with [1])
- [6] therefore r (combining [4] with [5])
- [7] The internal argument in steps [2–6] shows that $(p \wedge q) \rightarrow r$.

The above reasoning in [1–7] proves that $[p \rightarrow (q \rightarrow r)] \rightarrow [(p \wedge q) \rightarrow r]$.

(c) Same comment as for part (b). Note that here we have an internal argument and an “internal internal argument”.

- [1] Assumption: $(p \wedge q) \rightarrow r$
- [2] Sub-assumption: p
- [3] Sub-sub-assumption q
- [4] therefore $p \wedge q$ (combining [2] with [3])
- [5] therefore r (combining [4] with [1])
- [6] The internal internal argument in steps [3–5] shows that $q \rightarrow r$
- [7] The internal argument in steps [2–6] shows that $p \rightarrow (q \rightarrow r)$.

The above reasoning in [1–7] proves that $[(p \wedge q) \rightarrow r] \rightarrow [p \rightarrow (q \rightarrow r)]$.

(Wrapping up b and c) We have shown implications in both directions:

if $[p \rightarrow (q \rightarrow r)]$ then $[(p \wedge q) \rightarrow r]$, if $[(p \wedge q) \rightarrow r]$ then $[p \rightarrow (q \rightarrow r)]$.

Therefore $[p \rightarrow (q \rightarrow r)] \equiv [(p \wedge q) \rightarrow r]$.

Remark: There is another way to do (b) and (c) simultaneously, without doing everything we did above. Use the tautology $x \rightarrow y \equiv \neg x \vee y$ and De Morgan’s rules to show that $(p \wedge q) \rightarrow r \equiv \neg p \vee \neg q \vee r \equiv p \rightarrow (q \rightarrow r)$. The details are left to you.

4. (total 12 pts) Given the sequence a_1, a_2, a_3, \dots of real numbers, defined (recursively) by:

$$a_1 = \sqrt{2}, \quad \forall n \geq 1, \quad a_{n+1} = 1 + \frac{1}{a_n}.$$

a) (6 pts) Show that $\forall n \geq 1, \quad 1 < a_n < 2$.

b) (6 pts) Show that $\forall n \geq 1, \quad a_n$ is irrational. (You already know that $\sqrt{2}$ is irrational, so just prove the inductive step.)

Answer. (a) Let $P(n)$ be the statement $1 < a_n < 2$. We use induction to prove that $\forall n \geq 1, \quad P(n)$. First we prove the base case: $P(1)$ states that $1 < \sqrt{2} < 2$, which is true because we can deduce it by taking square roots of the inequalities $1 < 2 < 4$. Second, we prove the inductive step: let $n \geq 1$ be arbitrary, and assume $P(n)$. Thus $1 < a_n < 2$. Taking reciprocals shows that $1/2 < 1/a_n < 1$, and then adding 1 gives $3/2 < a_{n+1} = 1 + 1/a_n < 2$. This implies that $1 < a_{n+1} < 2$, so we obtain $P(n+1)$. The preceding discussion shows that $\forall n \geq 1, \quad P(n) \rightarrow P(n+1)$. Thus we have proved the inductive step. Now we combine the base case and the inductive step to conclude $\forall n \geq 1, \quad P(n)$ by induction. This completes the desired proof.

(b) Let $Q(n)$ be the statement that a_n is irrational. We already know $Q(1)$ from class, so we obtain the base case. The best way to prove the inductive step is by proving the **contrapositive** as follows. Let n be arbitrary and assume that a_{n+1} is **rational**. Then

there exist $p, q \in \mathbf{Z}$ such that $a_{n+1} = p/q$. But this means that $1 + 1/a_n = p/q$, in other words that $1/a_n = (p - q)/p$, and hence that $a_n = p/(p - q)$, with $p, p - q \in \mathbf{Z}$. This shows that a_n is rational. [Note that $p \neq q$ because we have already proved that the terms in the sequence are never equal to 1.] The above discussion shows that $\forall n \geq 1, (\neg Q(n+1)) \rightarrow (\neg Q(n))$. Replacing with the contrapositive shows that $\forall n \geq 1, Q(n) \rightarrow Q(n+1)$. This proves the inductive step. Combining the base case and the inductive step, we obtain by induction that $\forall n \geq 1, Q(n)$, which is what we wanted to prove.

5. (total 12 pts) Given sets X, Y , and a function $f : X \rightarrow Y$. Prove the following statements **FROM THE DEFINITION** of the image set and inverse image sets. I.e., you are **NOT ALLOWED** to use identities such as $f^{-1}(S \cap T) = f^{-1}(S) \cap f^{-1}(T)$, unless you include a proof of the identities you use. (It is possible to give a short solution that does not involve proving these identities first.)

a) (6 pts) Assume that $A, B, C \subseteq X$ satisfy $A \subseteq C$ and $B \subseteq C$. Show that $f(A) \cup f(B) \subseteq f(C)$.

b) (6 pts) Assume that $S, T \subseteq Y$ satisfy $S \cap T = \emptyset$. Show that $f^{-1}(S) \cap f^{-1}(T) = \emptyset$.

Answer. (a) We need to show an inclusion of sets. So we must start with an arbitrary element $y \in f(A) \cup f(B)$, and somehow conclude that y must belong to $f(C)$. This will show that $\forall y, y \in f(A) \cup f(B) \rightarrow y \in f(C)$, which will prove what we want. This is our plan for the proof.

We now carry out the plan: let $y \in f(A) \cup f(B)$ be arbitrary. We have **two cases**, (i) $y \in f(A)$ and (ii) $y \in f(B)$. If we are in case (i), then $\exists x \in A$ s.t. $y = f(x)$, by the definition of the set $f(A)$. But since $A \subseteq C$ and $x \in A$, we obtain that $x \in C$. Thus $y = f(x)$ for this specific $x \in C$, and this shows that $y \in f(C)$. On the other hand, if we are in case (ii), then $\exists x \in B$ s.t. $y = f(x)$. Once again, since $B \subseteq C$, we obtain $x \in C$ and $y = f(x) \in f(C)$. Thus in **both cases** we can conclude that $y \in f(C)$. Hence we have shown that $\forall y, y \in f(A) \cup f(B) \rightarrow y \in f(C)$, and this means that $f(A) \cup f(B) \subseteq f(C)$.

(b) This is easiest to prove by **contradiction**. Assume that $f^{-1}(S) \cap f^{-1}(T) \neq \emptyset$. Then $\exists x \in f^{-1}(S) \cap f^{-1}(T)$. In particular, $x \in f^{-1}(S)$, and so by the definition of $f^{-1}(S)$ we conclude that $f(x) \in S$. Similarly, $x \in f^{-1}(T)$ and this implies that $f(x) \in T$. Thus we have produced an element $f(x) \in S \cap T$, contradicting the given fact that $S \cap T = \emptyset$. The above reasoning by contradiction shows that $f^{-1}(S) \cap f^{-1}(T) = \emptyset$ after all.

6. (12 pts) Let $A = \{a, b, c, d\}$ and $B = \{1, 2, 3, 4\}$. How many functions $f : A \rightarrow B$ are there such that EITHER $[f(a) = 1]$ OR $[f(a) = 2 \text{ AND } f \text{ is injective}]$?

Answer. Define two sets $\mathcal{F} = \{\text{functions } f : A \rightarrow B \mid f(a) = 1\}$ and $\mathcal{G} = \{\text{functions } f : A \rightarrow B \mid (f(a) = 2) \wedge (f \text{ injective})\}$. Our task is to compute $|\mathcal{F} \cup \mathcal{G}|$. Note that $\mathcal{F} \cap \mathcal{G} = \emptyset$, since a given function cannot simultaneously satisfy $f(a) = 1$ and $f(a) = 2$ (whether or not f is injective). Thus $|\mathcal{F} \cup \mathcal{G}| = |\mathcal{F}| + |\mathcal{G}|$ by the addition principle.

We now compute $|\mathcal{F}|$ as follows: to choose $f \in \mathcal{F}$, we already know $f(a) = 1$, and we need to choose $f(b) \in \{1, 2, 3, 4\}$, followed by $f(c) \in \{1, 2, 3, 4\}$, followed by $f(d) \in \{1, 2, 3, 4\}$. By the multiplication principle, our total choices are $4 \cdot 4 \cdot 4 = 64$, and so $|\mathcal{F}| = 64$. [Note: this result can also be obtained by making a bijection between \mathcal{F} and the set of functions from $\{b, c, d\}$ to $\{1, 2, 3, 4\}$.]

We next compute $|\mathcal{G}|$. Once again, for $f \in \mathcal{G}$, we already know $f(a) = 2$. We must now choose $f(b) \in B - \{f(a)\} = \{1, 3, 4\}$ because $f(b)$ must be different from $f(a)$, so there are 3 choices for $f(b)$. We then choose $f(c) \in B - \{f(a), f(b)\}$, which gives us 2 choices for $f(c)$. Lastly we choose $f(d) \in B - \{f(a), f(b), f(c)\}$, which gives us 1 choice for $f(d)$. All in all we obtain $3 \cdot 2 \cdot 1 = 6$ choices for f , and so $|\mathcal{G}| = 6$. [Note: this result can also be obtained by making a bijection between \mathcal{G} and the set of injections from $\{b, c, d\}$ to $\{1, 3, 4\}$.]

Putting all of the above together, we obtain that $|\mathcal{F} \cup \mathcal{G}| = 64 + 6 = 70$.